

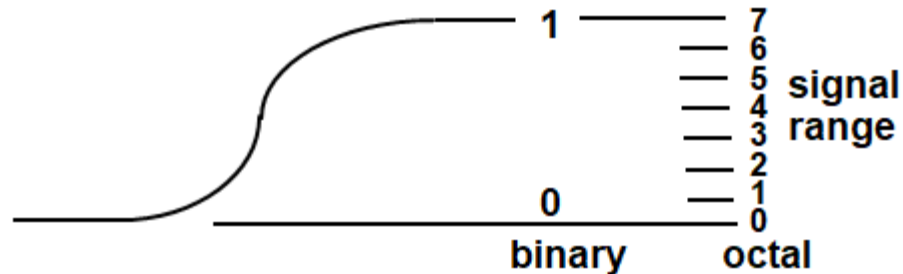
DIGITAL ELECTRONICS

LOGIC GATES

Digital Computers

- Imply that the computer deals with digital information, i.e., it deals with the information that is represented by binary digits
- Why *BINARY*? instead of Decimal or other number system ?

* Consider electronic signal



* Consider the calculation cost - Add

	0	1
0	0	1
1	1	10

	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18



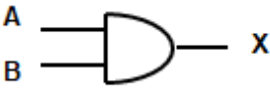







Types of Basic Logic Blocks

- **Combinational Logic Block**
Logic Blocks whose output logic value depends only on the input logic values
- **Sequential Logic Block**
Logic Blocks whose output logic value depends on the input values and the state (stored information) of the blocks

Functions of Gates can be described by

- Truth Table
- Boolean Function
- Karnaugh Map

COMBINATIONAL GATES

Name	Symbol	Function	Truth Table															
AND		$X = A \cdot B$ or $X = AB$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> </tr> <tr> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> </tr> </tbody> </table>	A	B	X	0	0	0	0	1	0	1	0	0	1	1	1
A	B	X																
0	0	0																
0	1	0																
1	0	0																
1	1	1																
OR		$X = A + B$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> </tr> </tbody> </table>	A	B	X	0	0	0	0	1	1	1	0	1	1	1	1
A	B	X																
0	0	0																
0	1	1																
1	0	1																
1	1	1																
I		$X = A'$	<table border="1"> <thead> <tr> <th>A</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	X	0	1	1	0									
A	X																	
0	1																	
1	0																	
Buffer		$X = A$	<table border="1"> <thead> <tr> <th>A</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> </tr> </tbody> </table>	A	X	0	0	1	1									
A	X																	
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NAND		$X = (AB)'$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	X	0	0	1	0	1	1	1	0	1	1	1	0
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NOR		$X = (A + B)'$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	X	0	0	1	0	1	0	1	0	0	1	1	0
A	B	X																
0	0	1																
0	1	0																
1	0	0																
1	1	0																
XOR Exclusive OR		$X = A \oplus B$ or $X = A'B + AB'$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>0</td> </tr> <tr> <td>0</td> <td>1</td> <td>1</td> </tr> <tr> <td>1</td> <td>0</td> <td>1</td> </tr> <tr> <td>1</td> <td>1</td> <td>0</td> </tr> </tbody> </table>	A	B	X	0	0	0	0	1	1	1	0	1	1	1	0
A	B	X																
0	0	0																
0	1	1																
1	0	1																
1	1	0																
XNOR Exclusive NOR or Equivalence		$X = (A \oplus B)'$ or $X = A'B' + AB$	<table border="1"> <thead> <tr> <th>A</th> <th>B</th> <th>X</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>1</td> </tr> <tr> <td>0</td> <td>1</td> <td>0</td> </tr> <tr> <td>1</td> <td>0</td> <td>0</td> </tr> <tr> <td>1</td> <td>1</td> <td>1</td> </tr> </tbody> </table>	A	B	X	0	0	1	0	1	0	1	0	0	1	1	1
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BOOLEAN ALGEBRA

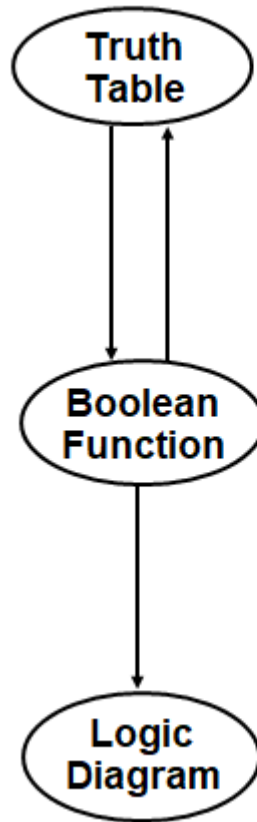
Boolean Algebra

- * Algebra with Binary(Boolean) Variable and Logic Operations
- * Boolean Algebra is useful in Analysis and Synthesis of Digital Logic Circuits
 - Input and Output signals can be represented by Boolean Variables, and
 - Function of the Digital Logic Circuits can be represented by Logic Operations, i.e., Boolean Function(s)
 - From a Boolean function, a logic diagram can be constructed using AND, OR, and I

Truth Table

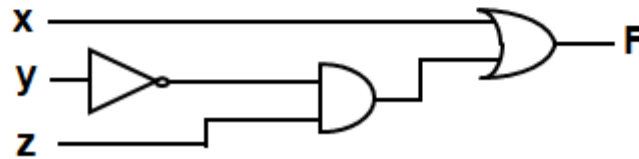
- * The most elementary specification of the function of a Digital Logic Circuit is the Truth Table
 - Table that describes the Output Values for all the combinations of the Input Values, called *MINTERMS*
 - n input variables $\rightarrow 2^n$ minterms

LOGIC CIRCUIT DESIGN



x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

$$F = x + y'z$$



COMPLEMENT OF FUNCTIONS

A Boolean function of a digital logic circuit is represented by only using logical variables and AND, OR, and Invert operators.

→ Complement of a Boolean function

- Replace all the variables and subexpressions in the parentheses appearing in the function expression with their respective complements

$$\begin{aligned} A, B, \dots, Z, a, b, \dots, z &\Rightarrow A', B', \dots, Z', a', b', \dots, z' \\ (p + q) &\Rightarrow (p + q)' \end{aligned}$$

- Replace all the operators with their respective complementary operators

$$\begin{aligned} \text{AND} &\Rightarrow \text{OR} \\ \text{OR} &\Rightarrow \text{AND} \end{aligned}$$

- Basically, extensive applications of the De Morgan's theorem

$$(x_1 + x_2 + \dots + x_n)' \Rightarrow x_1' x_2' \dots x_n'$$

$$(x_1 x_2 \dots x_n)' \Rightarrow x_1' + x_2' + \dots + x_n'$$

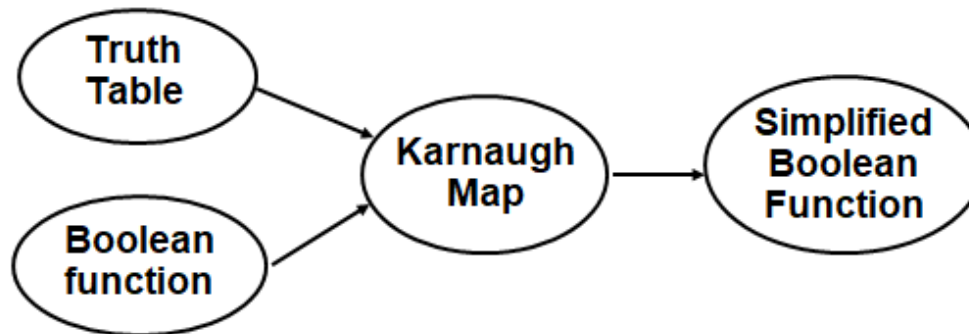
SIMPLIFICATION



Simplification from Boolean function

- Finding an equivalent expression that is least expensive to implement
- For a simple function, it is possible to obtain a simple expression for low cost implementation
- But, with complex functions, it is a very difficult task

Karnaugh Map (K-map) is a simple procedure for simplifying Boolean expressions.

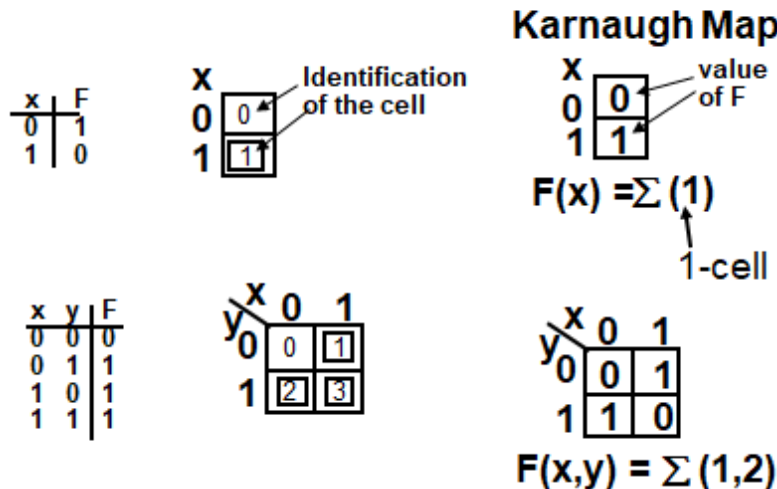


KARNAUGH MAP

Karnaugh Map for an n-input digital logic circuit (n-variable sum-of-products form of Boolean Function, or Truth Table) is

- Rectangle divided into 2^n cells
- Each cell is associated with a Minterm
- An output(function) value for each input value associated with a minterm is written in the cell representing the minterm
 → 1-cell, 0-cell

Each Minterm is identified by a decimal number whose binary representation is identical to the binary interpretation of the input values of the minterm.



KARNAUGH MAP

x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

		y			
		yz	00	01	11
x	0	0	1	3	2
	1	4	5	7	6

		yz			
		00	01	11	10
x	0	0	1	0	1
	1	1	0	0	0

$$F(x,y,z) = \Sigma (1,2,4)$$

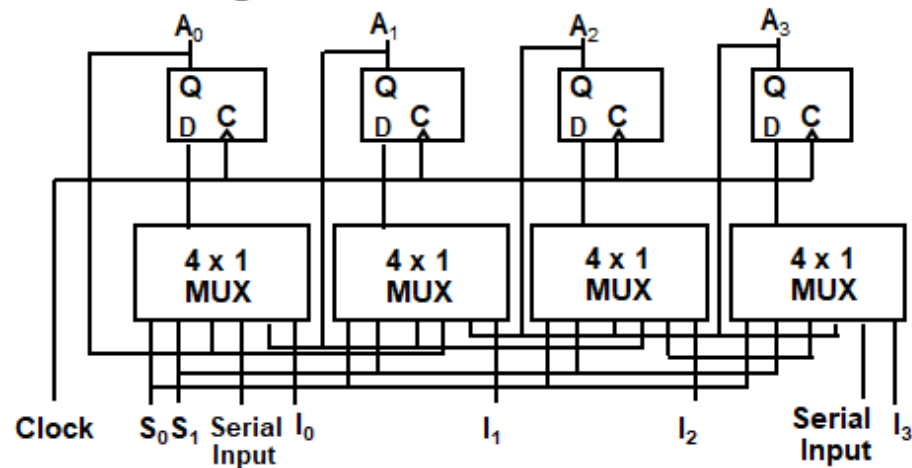
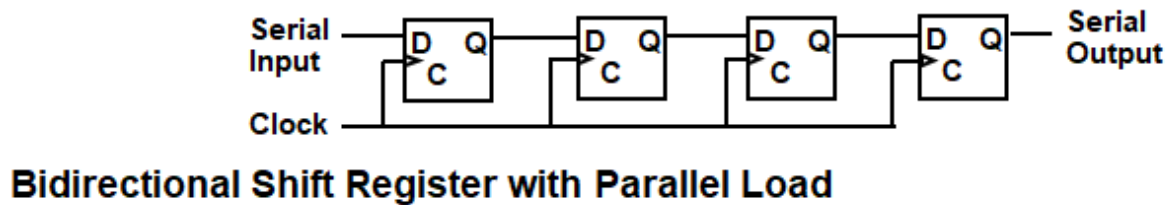
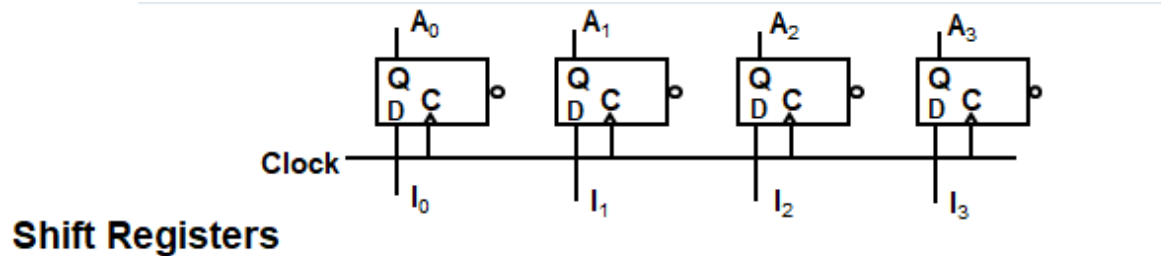
u	v	w	x	F
0	0	0	0	0
0	0	0	1	1
0	0	1	0	0
0	0	1	1	1
0	1	0	0	0
0	1	0	1	0
0	1	1	0	1
0	1	1	1	0
1	0	0	0	1
1	0	0	1	1
1	0	1	0	0
1	0	1	1	1
1	1	0	0	0
1	1	0	1	0
1	1	1	0	1
1	1	1	1	0

		w			
		wx	00	01	11
uv	00	0	1	3	2
	01	4	5	7	6
	11	12	13	15	14
	10	8	9	11	10

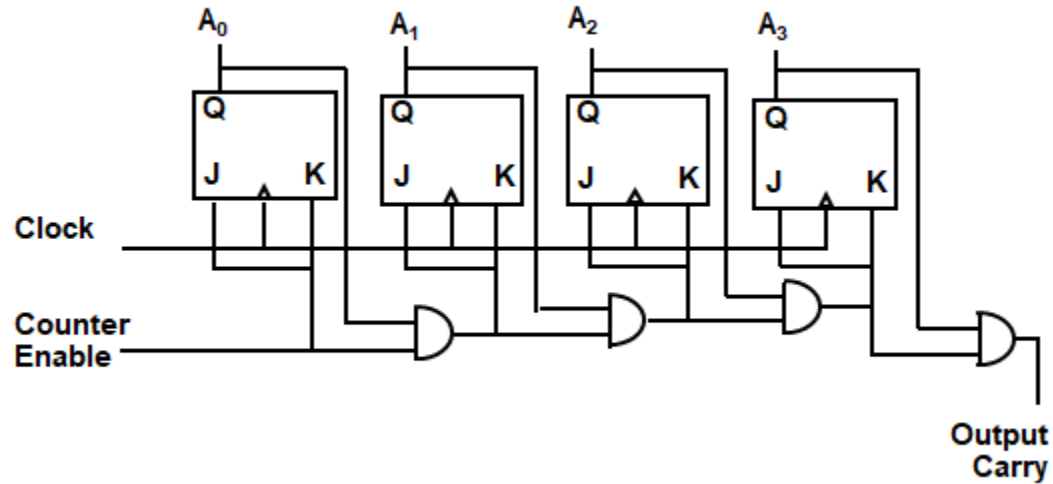
		wx			
		uv	00	01	11
uv	00	0	1	1	0
	01	0	0	0	1
	11	0	0	0	1
	10	1	1	1	0

$$F(u,v,w,x) = \Sigma (1,3,6,8,9,11,14)$$

SEQUENTIAL CIRCUITS - Registers

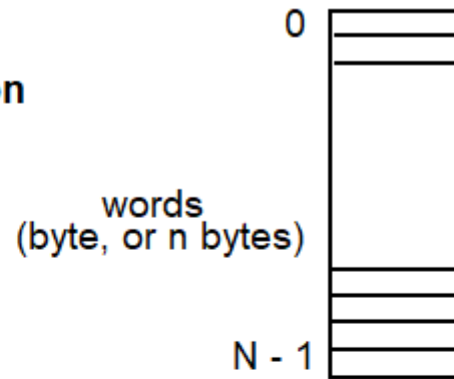


SEQUENTIAL CIRCUITS - Counters



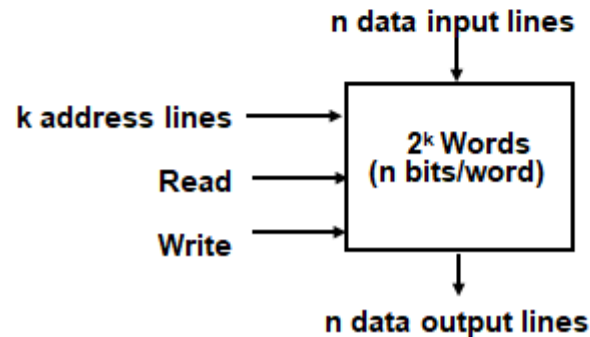
MEMORY COMPONENTS

Logical Organization



Random Access Memory

- Each word has a unique address
- Access to a word requires the same time independent of the location of the word
- Organization



TYPES OF ROM

ROM

- Store information (function) during production
- Mask is used in the production process
- Unalterable
- Low cost for large quantity production --> used in the final products

PROM (Programmable ROM)

- Store info electrically using PROM programmer at the user's site
- Unalterable
- Higher cost than ROM -> used in the system development phase
 - > Can be used in small quantity system

EPROM (Erasable PROM)

- Store info electrically using PROM programmer at the user's site
- Stored info is erasable (alterable) using UV light (electrically in some devices) and rewriteable
- Higher cost than PROM but reusable --> used in the system development phase. Not used in the system production due to eras ability

REFERENCES

- [www. Wikipedia.com](http://www.Wikipedia.com)
- www.slideshare.net